Indian Statistical Institute Semester-II 2012-2013 M.Tech.(CS) - First Year Class Test II (10 April, 2013) Subject: Automata, Languages and Computation Total: 20 marks To change an answer, scratch out the old answer and write the new answer clearly.

Do NOT overwrite.

Name:	_ Roll:
1. Fill in the blanks below with the correct expressions.	[1×5=5]
(a) Given a grammar $G = (V, T, P, S)$, a symbol $X \in$	$V \cup T$ is said to be
(i) useful if	;
(ii) nullable if	
(b) Given a grammar $G = (V, T, P, S)$, a production production if	on $A \to X_1 \dots X_n$ is said to be a unit
(c) Given a pushdown automaton (PDA) $P = (Q, Z)$	$\Sigma, \Gamma, \delta, q_0, Z_0, F$), the language accepted

- by P(i) by **final state** is given by $L(P) = \{w \mid \vdash_P^* \};$
 - (ii) by **empty stack** is given by $N(P) = \{w \mid \vdash_P^* \}$.
- Recall that *postfix notation* is a method for writing arithmetic expressions in which every operator is written after all of its operands. For example, the postfix equivalent of A × B + C/D is AB×CD/+. Write a context-free grammar (CFG) for arithmetic expressions in postfix notation involving variables and the operators +, -, × and /. You may assume that variable names consist of single letters only (as in the example above). You should use a single non-terminal S. [3] Answer:

3. Let L = {w | w is obtained by taking a syntactically correct C program and removing everything other than the keywords if and else from it}. Draw the state diagram of a PDA that accepts L by empty stack. You may assume that if and else are single symbols. [7]

4. Let G = (V, T, P, S) be a CFG in Greibach Normal Form. Let |V| = n, |T| = m, |P| = p. Suppose that p_0 of the productions are of the form $A_0 \rightarrow aA_1A_2...A_k$ where k is a fixed number, $A_i \in V$ for $0 \le i \le k$ and $a \in T$. The remaining productions are of the form $A \rightarrow a$ where $a \in T$. Let G' = (V', T, P', S) be a CFG in Chomsky Normal Form (CNF) obtained from G using the standard algorithm for conversion to CNF. Then:

 $|V'| \leq \underline{\qquad} |P'| \leq \underline{\qquad}$

Your bounds should be tight. Briefly justify your answer.

[5]